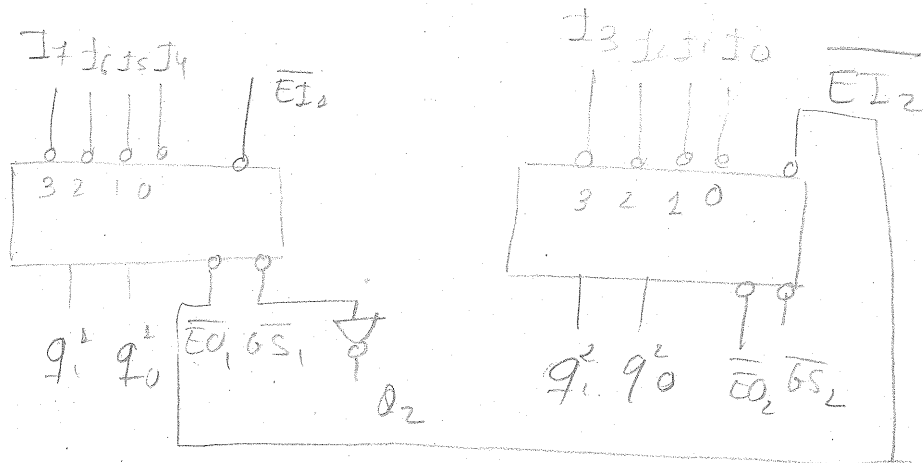


Problemas



$Q_2 = 1$ si I_7, I_6, I_5, I_4 están activas

$Q_2 = 0$ si I_3, I_2, I_1, I_0 están activas
o no hay ninguna activa

$$Q_2 = \overline{GS}$$

$Q_1 = 1$ si $q_1 = 1$
o $q_0 = 1$

$$Q_2 = Q_1 Q_0$$

1	1	1
1	1	0
0	1	1
0	1	0

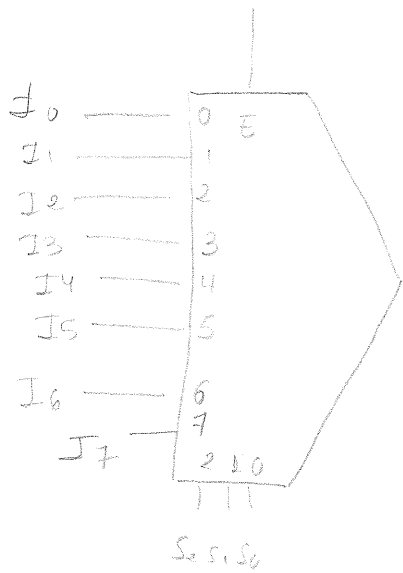
$$Q_1 = q_1 + q_0$$

$$Q_0 = q_1 + q_0$$

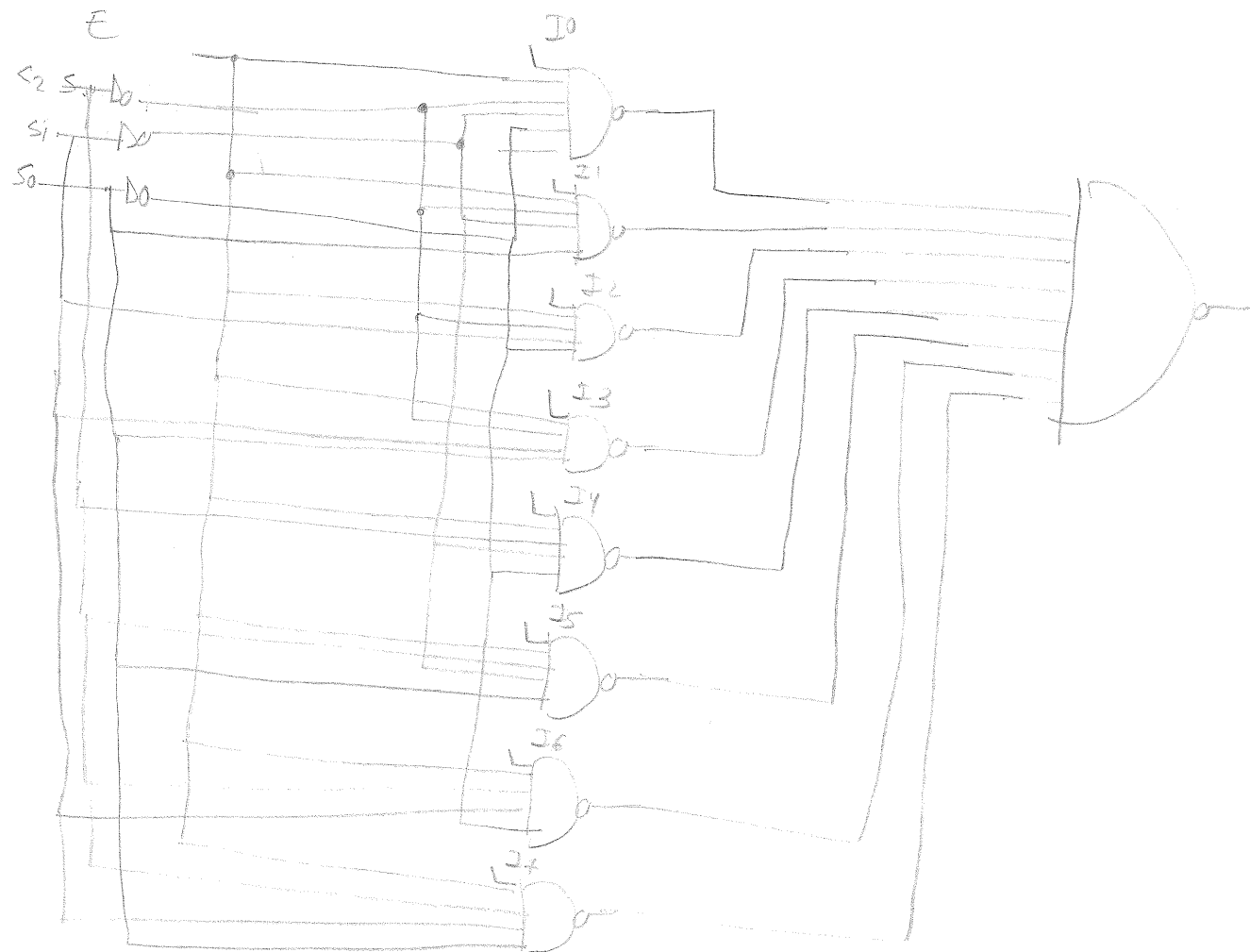
$$\overline{E_0} = \overline{E_0} Q_2$$

$$\overline{GS} = \overline{GS_1} \cdot \overline{GS_2}$$

PROBLEMA 8



$$O = E (\bar{S}_2 \bar{S}_1 \bar{S}_0 I_0 + \bar{S}_2 \bar{S}_1 S_0 I_1 + \bar{S}_2 S_1 \bar{S}_0 I_2 + \dots + S_2 S_1 S_0 I_7)$$

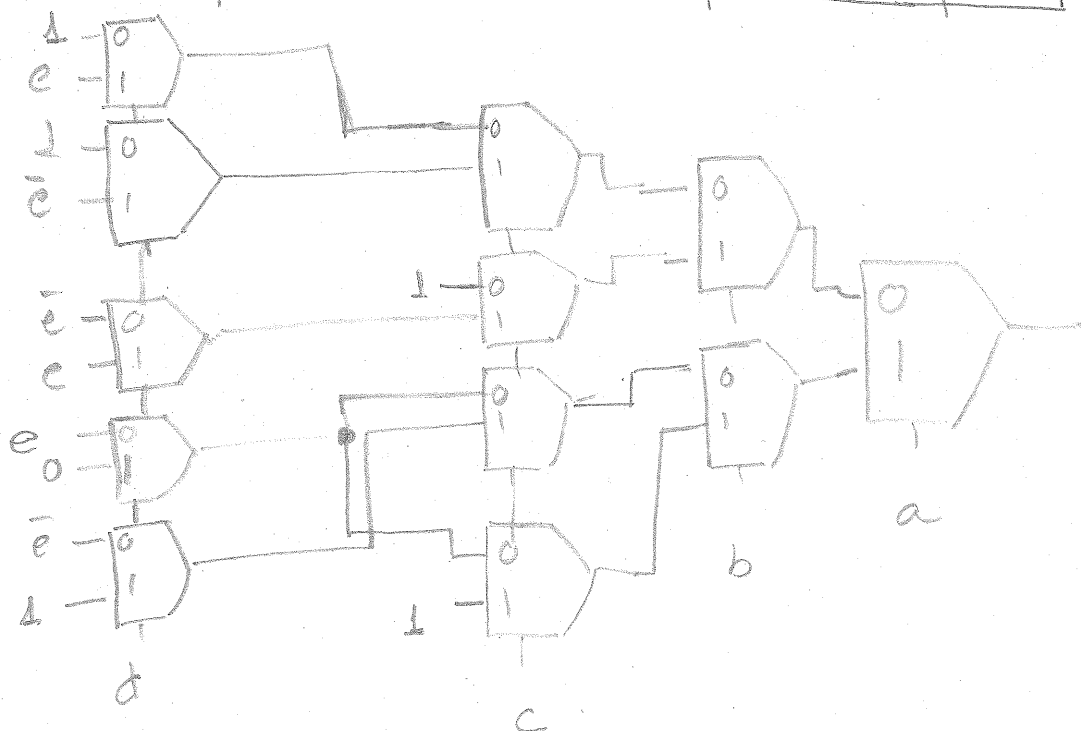


Solution P. 92 . Bol. 4

Realize com MUX 2:1 a função

$$F = \Sigma (0, 1, 3, 4, 5, 6, 8, 9, 10, 11, 12, 15, 17, 20, 22, 23, 25, 26, 29, 30, 31)$$

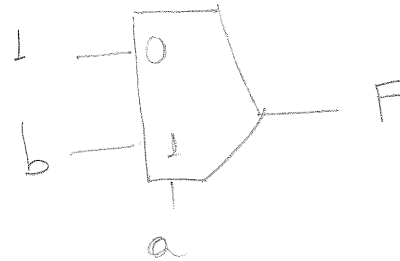
		abc							
de		000	001	011	010	110	111	101	100
		1	1	1	1		1	1	
01		1	1		1	1	1		1
11		1		1	1		1	1	
10			1		1		1	1	



PROBLEMA 10

$$F(a,b,c) = \sum (0, 3, 7) + d(1, 2, 6)$$

c \ ab	00	01	11	10
0	1	d	d	
1	d	1	1	



$$F(a,b,c,d,e) = \sum (2,3,4,5,6,7,8,9,10,14,15,16,17,18,19,20,21)$$

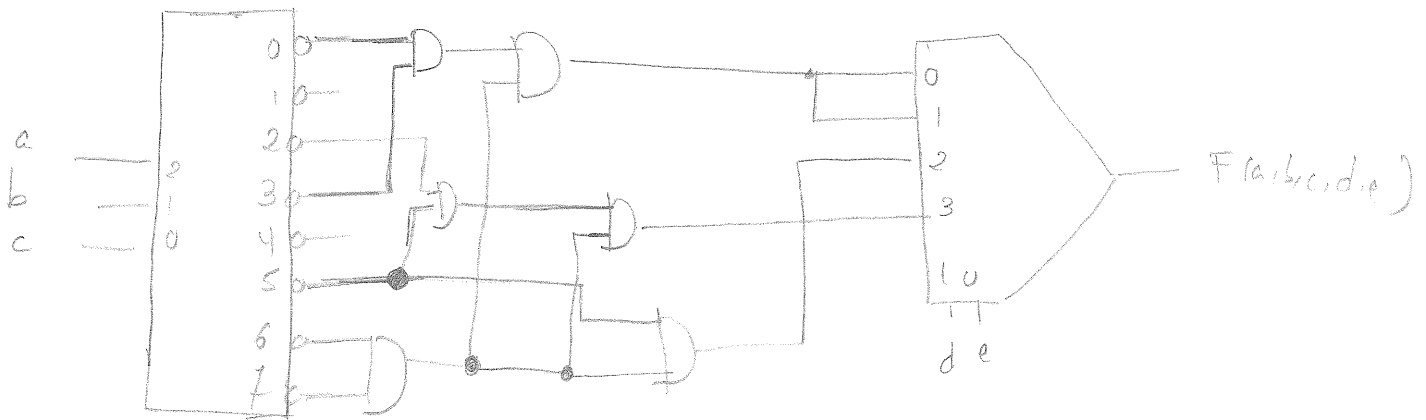
Un único MUX 4:1

Un único DEC 3:8

Puertas AND de dos entradas

Variables en único rail.

de	abc							
	000	001	011	010	110	111	101	100
00		1		1			1	1
01		1		1			1	1
11	1	1	1					1
10	1	1	1	1				1



Problema 12

$$F = \sum (1, 3, 11, 13, 21, 23, 25, 31) + d(5, 19, 27)$$

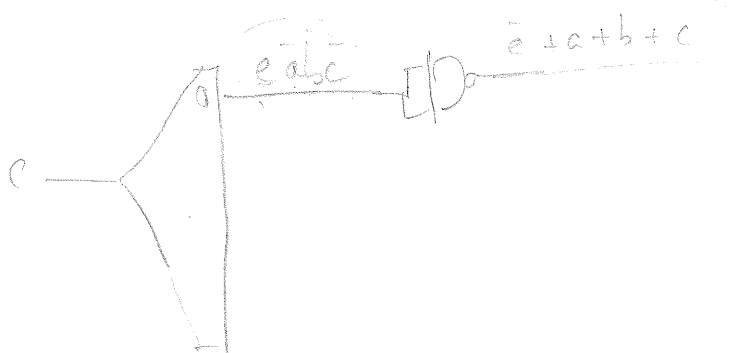
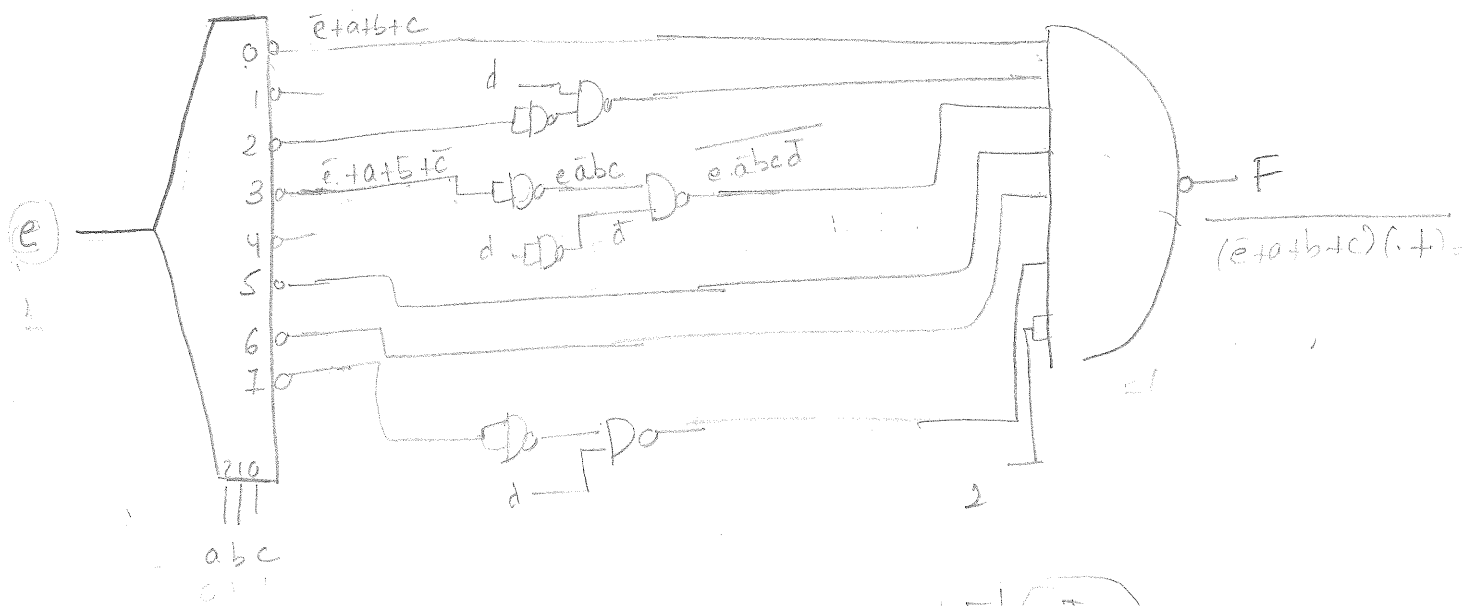
1 decimo DEMUX 1: 8

1 NAND: 7 entradas

Varus NAND de dos entradas

		abc							
		000	001	011	010	110	111	101	100
de	00								1
	01	1	d	1		1		1	
	11	1			1	d	1	1	d
	10								

F



$$\bar{a} \bar{b} \bar{c} \bar{d} \bar{e}$$

$$a + b + \bar{c} + d + \bar{e}$$

Problema 13

P. 1

$$F = \sum (0, 1, 3, 7, 9, 12, 15)$$

$$G = \prod (0, 1, 2, 5, 6, 10, 11)$$

$$H = (X_3 + \bar{X}_2) \cdot (X_2 + \bar{X}_1 + X_0) = \sum (0, 1, 3, 8, 9, 11, 12, 13, 14, 15)$$

$X_1 X_0$	$X_3 X_2$			
	00	01	11	10
00	1	0	1	1
01	1	0	1	1
11	1	0	1	1
10	0	0	1	0

ROM $2^4:3$

POS	CONT (F&H)
\$0	5
\$1	5
\$2	0
\$3	7
\$4	2
\$5	0
\$6	0
\$7	6
\$8	3
\$9	7
\$A	0
\$B	1
\$C	7
\$D	3
\$E	3
\$F	7

Tabla de programación de la ROM

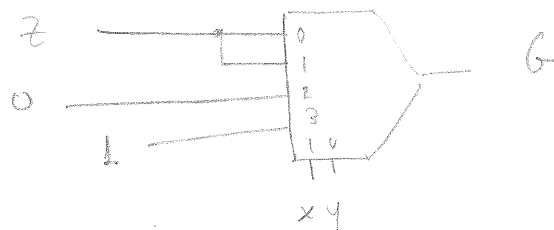
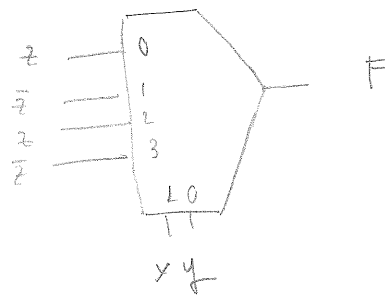
Pos	CONT
0	0 0
1	1 0
2	1 0
3	0 1
4	1 0
5	0 1
6	0 1
7	1 1
$x y z$	F G

z	xy	00	01	11	10
0		0	1	0	1
1		1	0	1	0

$$F = x \oplus y \oplus z$$

z	xy	00	01	11	10
0				1	
1		1	1	1	

$$G = \bar{x}z + xy$$



$$F = \sum (10, 1, 3, 7, 9, 12, 15)$$

$$PLA(4, 8, 3)$$

$$G = \prod (0, 1, 2, 5, 6, 10, 11)$$

$$H = (x_3 + \bar{x}_2)(x_2 + \bar{x}_1 + x_0)$$

$x_3 x_2$	00	01	11	10
$x_1 x_0$	1	0	1	0
01	1	0	0	1
11	1	1	1	0
10	0	0	0	0

$x_3 x_2$	00	01	11	10
$x_1 x_0$	0	1	1	1
01	0	0	1	1
11	1	1	1	0
10	0	0	1	0

$x_3 x_2$	00	01	11	10
$x_1 x_0$	1	0	1	1
01	1	0	1	1
11	1	0	1	0
10	0	0	1	0

$$F = p_1 + p_2 + p_6 + p_7 + p_8$$

$$G = p_2 + p_3 + p_4$$

$$H = p_1 + p_3 + p_4 + p_5$$

$$p_1 = \bar{x}_3 \bar{x}_2 \bar{x}_1$$

$$p_2 = \bar{x}_3 x_1 x_0$$

$$p_3 = x_3 x_2$$

$$p_4 = x_3 \bar{x}_1$$

$$p_5 = \bar{x}_2 x_1 x_0$$

$$p_6 = x_2 x_1 x_0$$

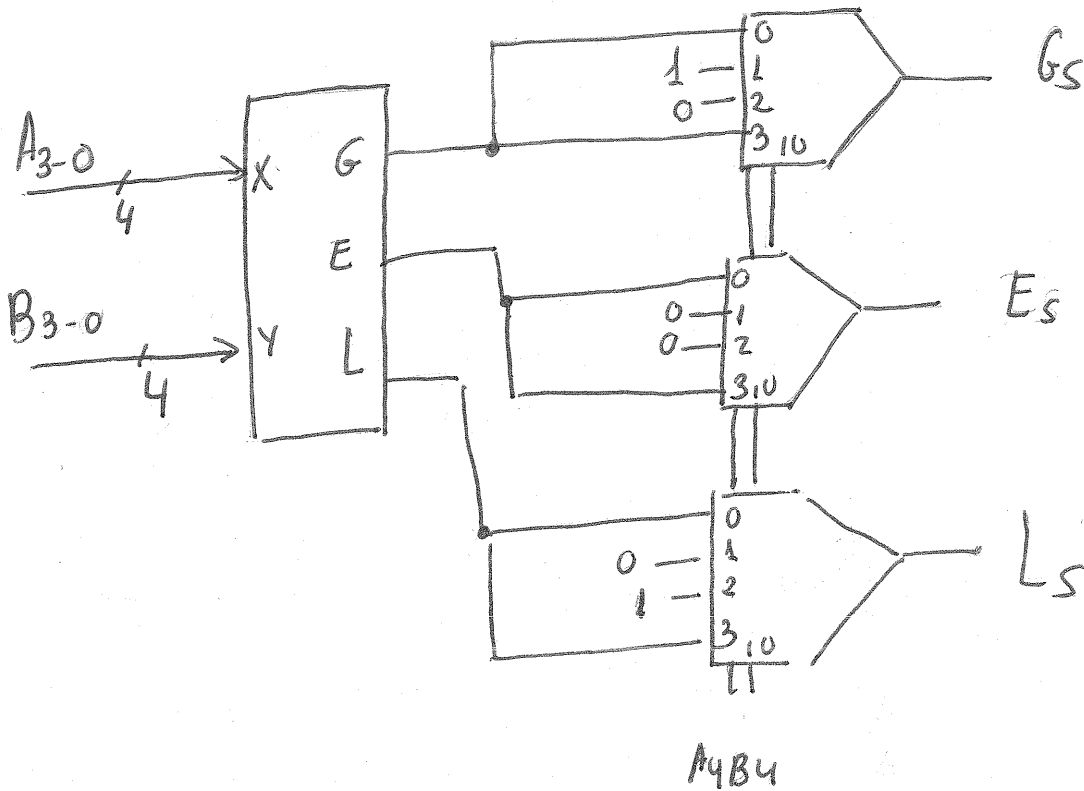
$$p_7 = x_3 x_2 \bar{x}_3 \bar{x}_0$$

$$p_8 = \bar{x}_2 \bar{x}_1 x_0$$

PROBLEMA 17

$\boxed{10000} \rightarrow -15$ el de menor magnitud
 $\boxed{10101} \rightarrow -10$ es el más pequeño

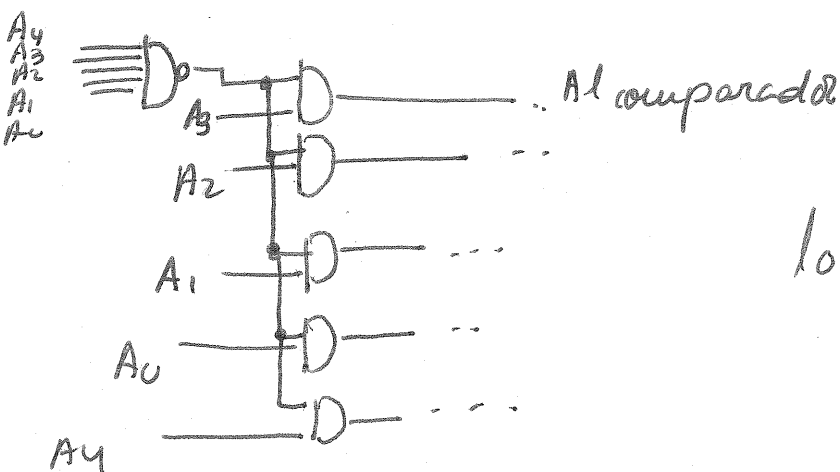
(a)



(b)

$\begin{matrix} 11111 \\ 00000 \end{matrix} \left\{ \begin{array}{l} \text{Se tendría que activar } E \\ \text{pero no va a ser así} \end{array} \right.$

Es la única situación en la que se produce un problema. Hay que transformar el 11111 en 00000



PROBLEMA 18

(a) $f = x_1 x_2 + \bar{x}_1 (\bar{x}_2 x_3 + x_2) =$

$$= x_1 x_2 + \bar{x}_1 (x_3 + x_2) = x_1 x_2 + \bar{x}_1 x_3 + \bar{x}_1 x_2 = x_2 + \bar{x}_1 x_3$$

(b)

$$f = \bar{x}_1 \bar{x}_3 (\bar{x}_4 \bar{x}_5 + \bar{x}_4 x_5 x_2) + \bar{x}_1 x_3 (\bar{x}_4 \bar{x}_5 + \bar{x}_4 x_5 \bar{x}_2 + x_4 x_5) +$$

$$+ x_1 \bar{x}_3 (x_4 \bar{x}_5 x_2) + x_1 x_3 (\bar{x}_4 x_5 x_2 + x_4 \bar{x}_5 x_2) =$$

$$= \bar{x}_1 \bar{x}_3 \bar{x}_4 \bar{x}_5 + \bar{x}_1 \bar{x}_3 \bar{x}_4 x_5 x_2 + \bar{x}_1 x_3 \bar{x}_4 \bar{x}_5 + \bar{x}_1 x_3 \bar{x}_4 x_5 \bar{x}_2 +$$

$$+ \bar{x}_1 x_3 x_4 x_5 + x_1 \bar{x}_3 x_4 \bar{x}_5 x_2 + x_1 x_3 \bar{x}_4 x_5 x_2 + x_1 x_3 x_4 \bar{x}_5 x_2$$

$x_4 x_5$	$x_1 x_2 x_3$							
	000	001	011	010	110	111	101	100
00	1	1	1	1				
01		1		1		1		
11		1	1					
10			1			1		

$$f = \bar{x}_1 \bar{x}_4 \bar{x}_5 + \bar{x}_1 \bar{x}_2 x_3 x_5 + \bar{x}_1 x_3 x_4 x_5 + \bar{x}_1 x_2 \bar{x}_3 \bar{x}_4 + x_2 x_3 x_4 \bar{x}_5 +$$

$$+ x_1 x_2 x_3 \bar{x}_4 x_5$$

PROBLEMA 18

(c)

$$f = E(\bar{x}(\bar{y}+z) + x(\bar{y}+\bar{z}))$$

$$E = \overline{(y+z)(\bar{y}+\bar{z})} = \bar{z}$$

$$f = \bar{z}(\bar{x}\bar{y} + \bar{x}z + x\bar{y} + x\bar{z}) = \bar{z}(\bar{y} + \bar{x}z + x\bar{z}) = \bar{z}\bar{y} + x\bar{z}$$

PROBLEMA 19

$$F = \sum (1, 2, 3, 4, 6, 7, 8, 9, 14)$$

Al obtenerse una s.p. negada, la función que tengo que programar en la PAL es $\bar{F} = \sum (0, 5, 10, 11, 12, 13, 15)$

cd \ ab				
	00	01	11	10
00	1		1	
01		1	1	
11			1	1
10				1

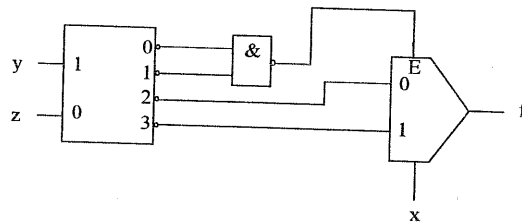
$$\bar{F} = \underbrace{\bar{a}\bar{b}\bar{c}\bar{d} + b\bar{c}d + abc}_{\bar{f}} + acd + a\bar{b}c$$

$$\bar{F} = \bar{f} + acd + a\bar{b}c$$

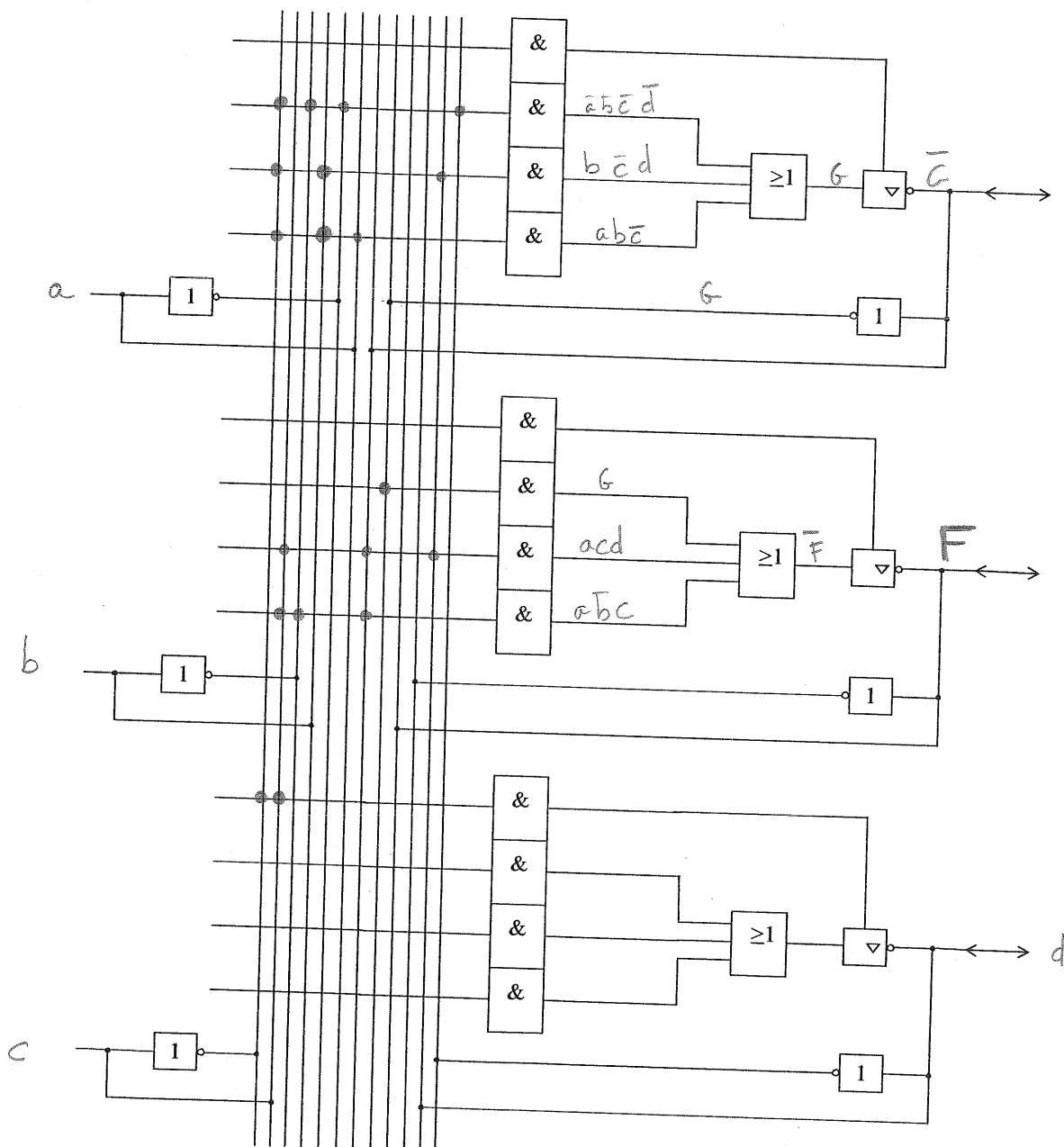
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c)



P19. Realice la función $F = \Sigma (1, 2, 3, 4, 6, 7, 8, 9, 14)$, mediante la PAL de la figura .

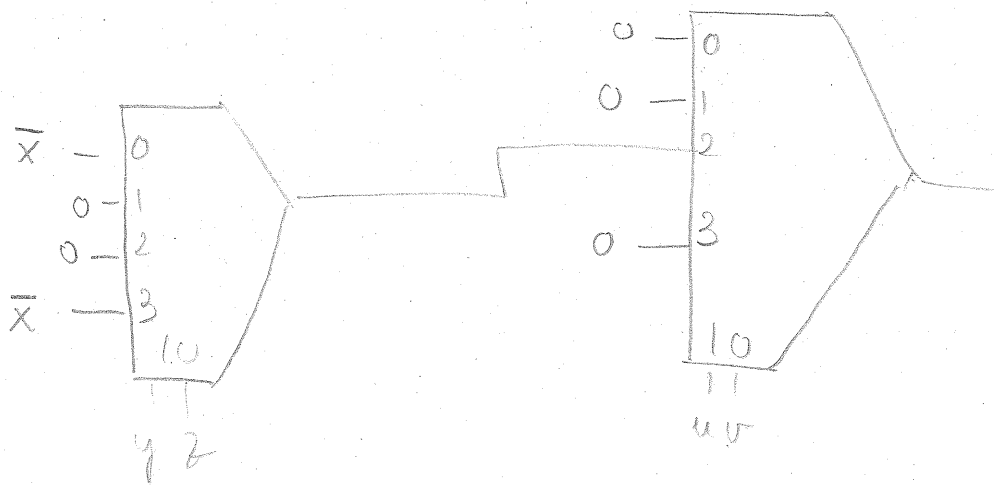


PROBLEMA 20

(a)

$x y z$		000	001	011	010	110	111	101	100
uv	00	0	0	0	0	0	0	0	0
	01	0	0	0	0	0	0	0	0
	11	0	0	0	0	0	0	0	0
	10	1	0	1	0	0	0	0	0

(b)



PROBLEMA 21

$$f = \bar{x} D_3 + x \bar{y} + \bar{x} D_2 + x \bar{y} + \bar{x} D_1 + x \bar{y} + \bar{x} D_0 =$$

$$= \bar{x} D_3 + x + \bar{x} D_2 + \bar{x} D_1 + \bar{x} D_0 =$$

$$x + \bar{x} (D_3 + D_2 + D_1 + D_0)$$

De la ROM sólo se utilizan las posiciones
2 y 3

POS	CONT			
	D_3	D_2	D_1	D_0
010	1	0	0	0
011	0	0	0	0

Por lo tanto

$$D_2 = D_1 = D_0 = 0$$

$$D_3 = \bar{z}$$

$$f = x + \bar{x} \bar{z} = x + \bar{z}$$

PROBLEMA 22

El sistema permite la visualización en los displays de 7 segmentos del contenido de la ROM

PROBLEMA 23

$$f = 1 \quad \text{si} \quad x_1 x_0 \geq y_1 y_0$$

PROBLEMA 24

$$1 = \overline{D_1} \overline{D_0} D_3 + \overline{D_1} D_0 D_3 + D_1 \overline{D_0} D_3 = \overline{D_1} D_3 + D_1 \overline{D_0} D_3 =$$

$$= P_3 (\overline{D_1} + D_1 \overline{D_0}) = \overline{D_3} + P_1 D_0$$

$$f_2 = D_1 P_0 D_3$$

Debemos calcular D_1 , D_0 y D_3 en función de las variables de entrada para ello nos basamos en la programación de la ROM.

POS			CONT			
b	a	c	D ₃	D ₂	D ₁	D ₀
0	0	0	1	0	1	0
0	0	1	1	1	0	1
0	1	0	0	0	1	0
0	1	1	1	0	1	1
1	0	0	1	1	0	0
1	0	1	0	1	1	1
1	1	0	0	0	1	1
1	1	1	0	1	1	1

	ba	00	01	11	10
c					
0	1				1
1	1	1			

$$D_3 = \overline{b}c + \overline{a}\overline{c}$$

ba

	00	01	11	10
0	1	1	1	
1		1	1	1

$$b_1 = \bar{b}\bar{c} + a + bc$$

ba

	ol	ll	lb
o		1	
l	1	1	1

$$b \circ = c + ba$$

PROBLEMA 24

Sustituyendo en las funciones:

$$f_1 = \overline{(b\bar{c} + a\bar{c})} + (b\bar{c} + a + bc)(c + ba) =$$

$$= (b + \bar{c})(a + c) + ac + ab + bc = ab + a\bar{c} + ac + ab + bc = a + bc$$

$$f_2 = (b\bar{c} + a + bc)(c + ba)(b\bar{c} + a\bar{c}) =$$

$$= (ac + ab + bc)(b\bar{c} + a\bar{c}) = a\bar{b}c$$

Problema 25

Si $x_3 x_2 x_1 x_0 = y_3 y_2 y_1 y_0$, $z_3 z_2 z_1 z_0 = 0000$

Si $x_3 x_2 x_1 x_0 > y_3 y_2 y_1 y_0$, $z_3 z_2 z_1 z_0 = x_3 x_2 x_1 x_0$

Y a fin en este caso $A < B = L$

Si $x_3 x_2 x_1 x_0 < y_3 y_2 y_1 y_0$, $z_3 z_2 z_1 z_0 = y_3 y_2 y_1 y_0$

Y a fin en este caso $A < B = 0$

- A la salida se obtiene el mayor de los dos números de entrada o $\emptyset$ si son iguales

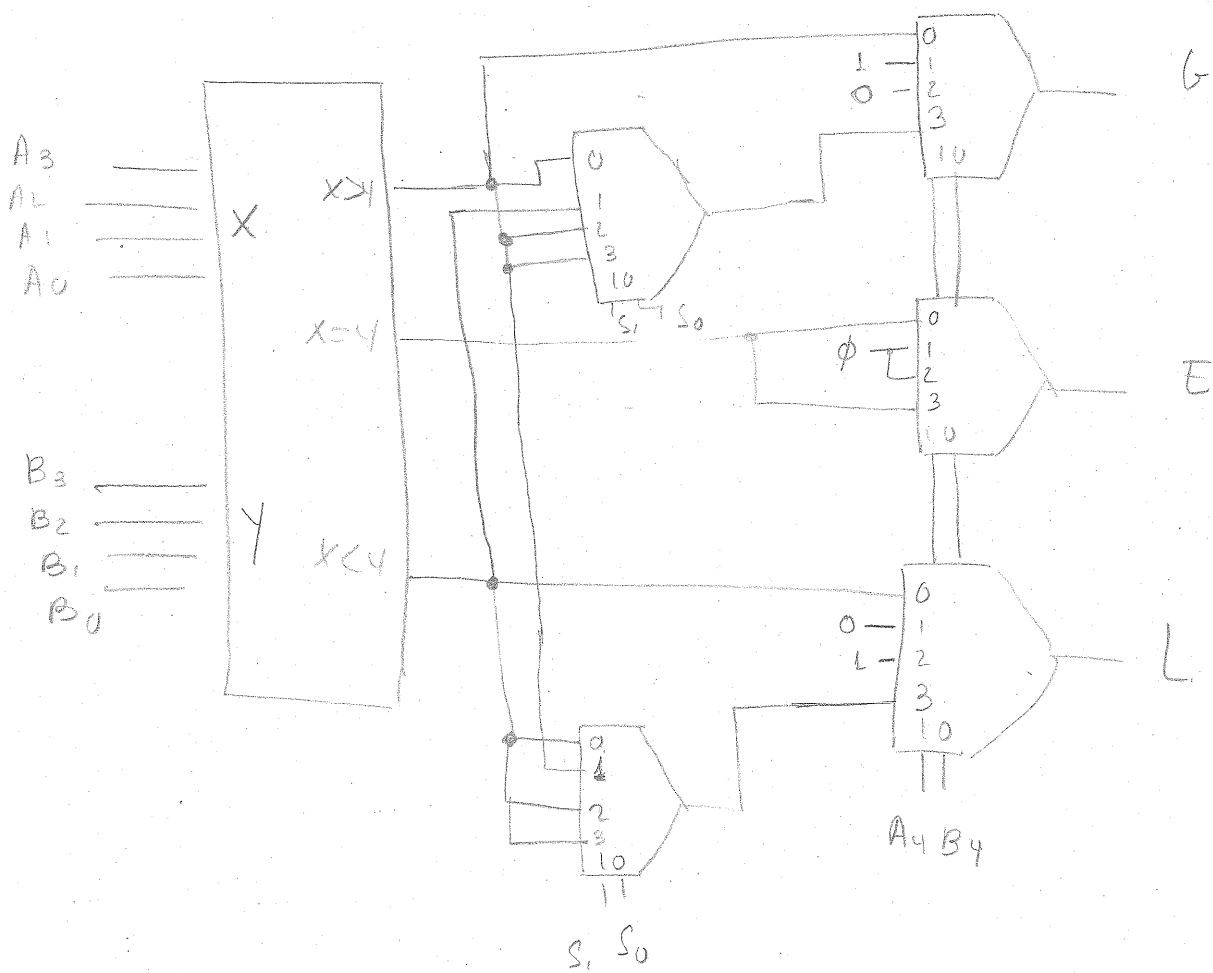
El circuito puede diseñarse con una ROM

$2^8 \times 4$.

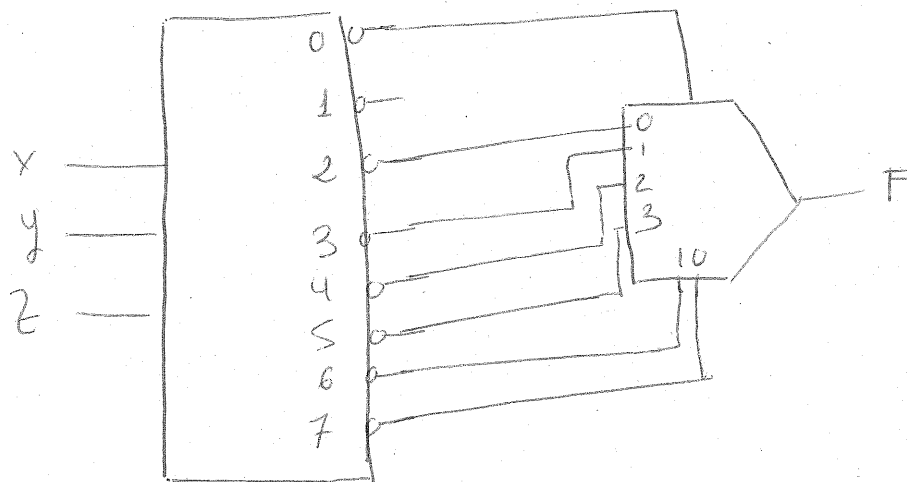
A continuación se indica el contenido de la ROM para los casos mencionados en el enunciado

POS					CONT				
$x_3 x_2 x_1 x_0$	y_3	y_2	y_1	y_0	z_3	z_2	z_1	z_0	
\$ 1 0					\$ 1				
\$ 1 1					\$ 0				
\$ 1 2					\$ 2				
\$ 6 7					\$ 7				
\$ 8 4					\$ 8				
\$ A A					\$ 0				
\$ D F					\$ F				

Problema 26

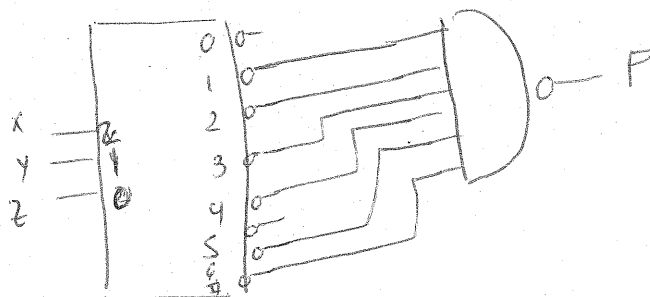


PROBLEMA 44



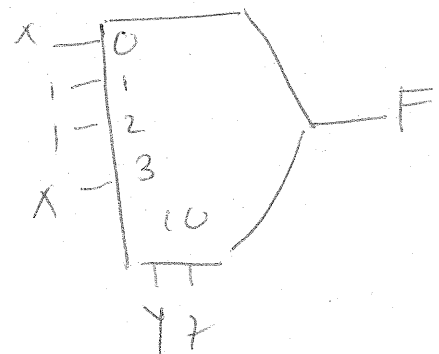
$$F = \prod (0, 5)$$

(a.)



(s)

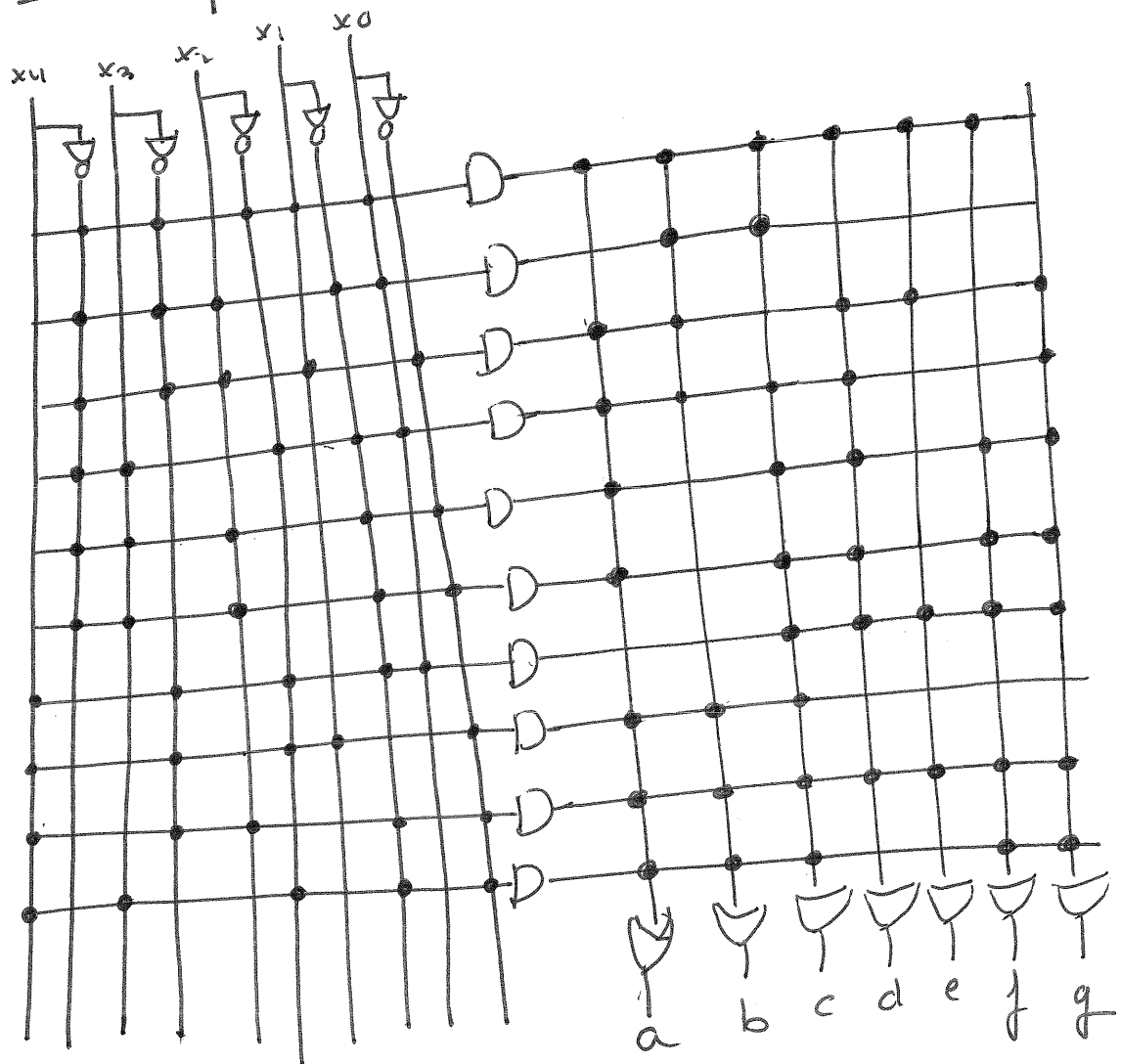
z	xy			
	00	01	11	10
0	0	1	1	1
1	1	1	1	0



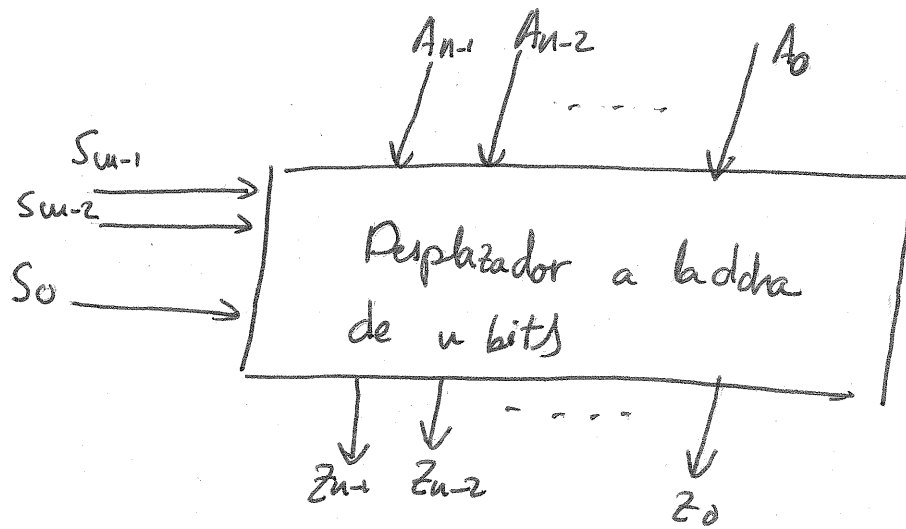
PROBLEMA 46

x_4	x_3	x_2	x_1	x_0	a	b	c	d	e	f	g
0	0	0	1	1	1	1	1	1	1	1	0
0	0	1	0	1	0	1	1	0	0	0	0
0	0	1	1	0	1	1	0	1	1	0	1
0	1	0	0	1	1	1	1	0	0	1	1
0	1	0	1	0	1	1	1	0	0	1	1
0	1	1	0	0	1	0	1	1	0	1	1
1	0	0	0	1	0	0	1	1	1	1	1
1	0	0	1	0	1	1	1	0	0	0	0
1	0	1	0	0	1	1	1	1	1	1	1
1	1	0	0	0	1	1	0	0	1	1	1

$\begin{array}{c|c|c} a & & b \\ \hline f & g & \\ e & & c \\ \hline a & & \end{array}$

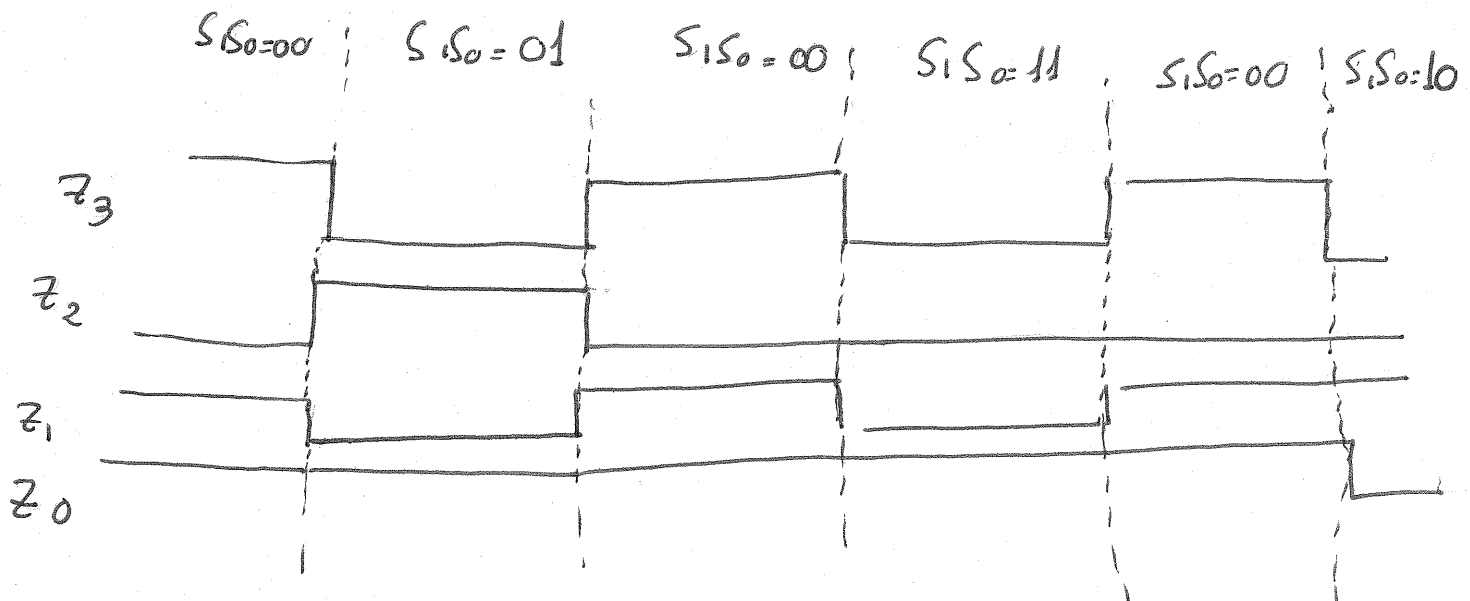
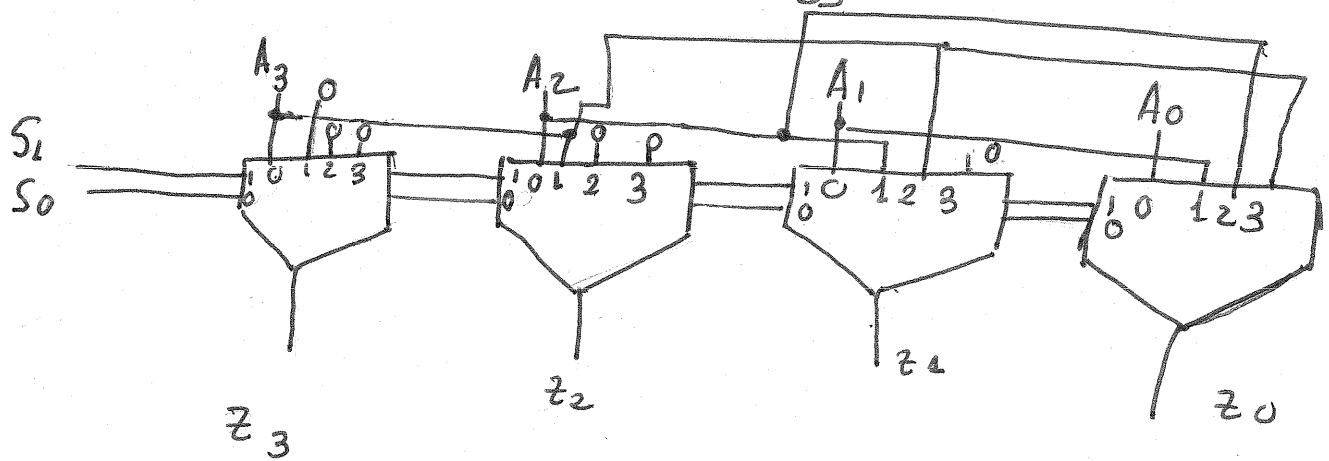


PROBLEMA 52



(a) $n=4$
 $m=2$

4 MUX's de 4 canales



c) dividir entre las potencias de dos